CS 591 K1: **Data Stream Processing and Analytics** Spring 2021 Window aggregation

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Window operators



Practical way to perform operations on unbounded input

• e.g. joins, holistic aggregates

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 - e.g. joins, holistic aggregates
- Compute on most *recent* events only
 - happened 2 hours ago

Window operators

• when providing real-time traffic information, you probably don't care about an accident that



- Practical way to perform operations on unbounded input
 - e.g. joins, holistic aggregates
- Compute on most *recent* events only
 - ullethappened 2 hours ago
- Recent might mean different things
 - last 5 sec ullet
 - last 10 events
 - last 1h every 10 min
 - last user session lacksquare

Window operators

when providing real-time traffic information, you probably don't care about an accident that



Keyed vs. non-keyed windows

Window operators can be applied on a keyed or a non-keyed stream:

- Window operators on keyed windows are evaluated in parallel
- Non-keyed windows are processed in a single thread

- A window assigner determines how the elements of the input stream are grouped into windows.
- A window function is applied on the window contents and processes the elements assigned to each window.

To create a window operator, you need to specify two window components:



Window types

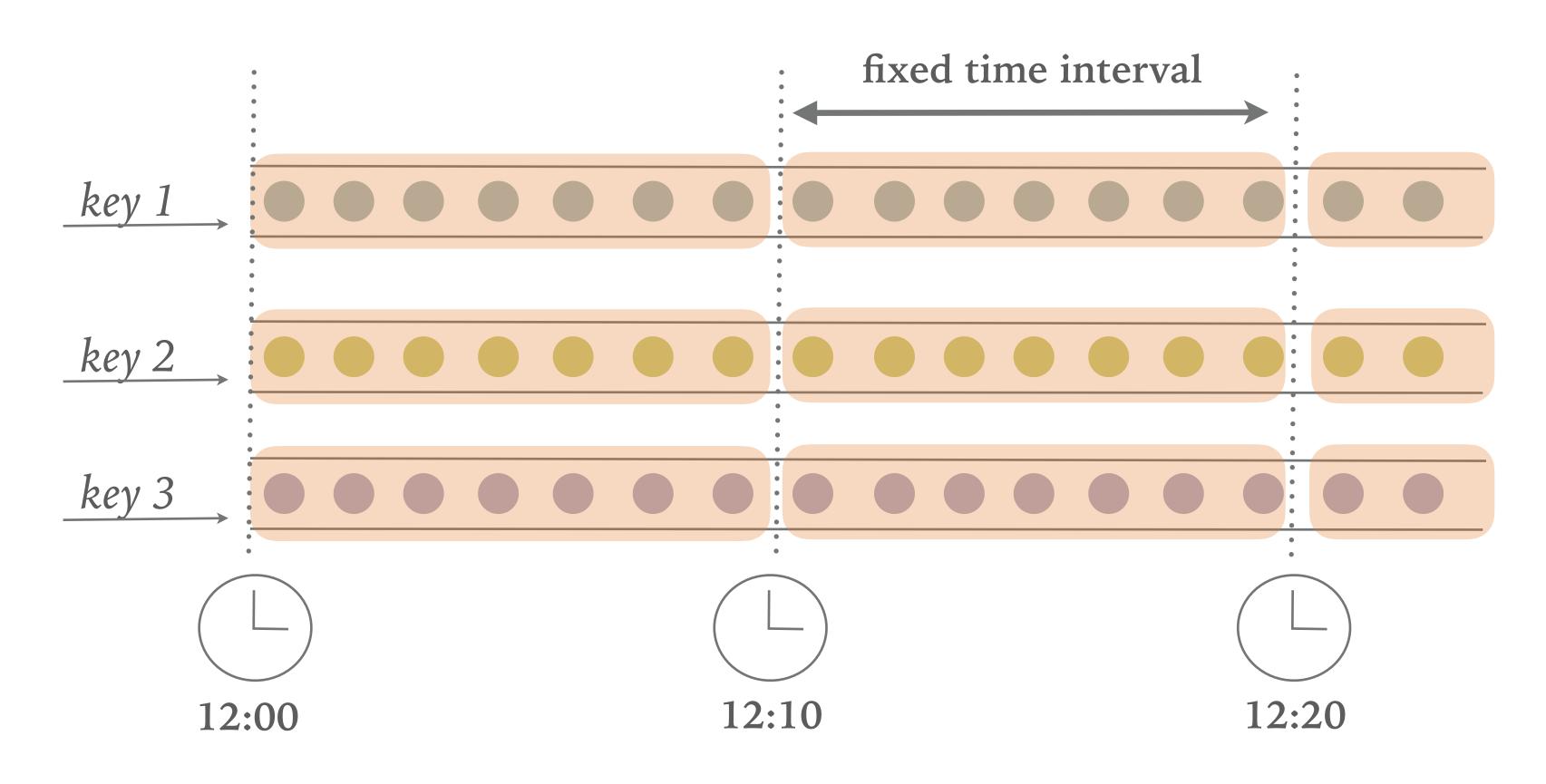
a continuous stream of events:

- **Tumbling** (fixed) windows split time into segments of equal length 1. The end of a window marks the start of the next one.
 - Each event belongs to one window only.
- Sliding (hopping) windows further define a slide parameter Is which determines how often a new window starts. Consecutive windows overlap when $l_s < l_s$
 - Events may belong to multiple windows.
- Session windows define a period of activity followed by a period of inactivity. A session window ends if no event arrives for some time gap l_g .

The window type defines the logic based on which we derive finite windows from

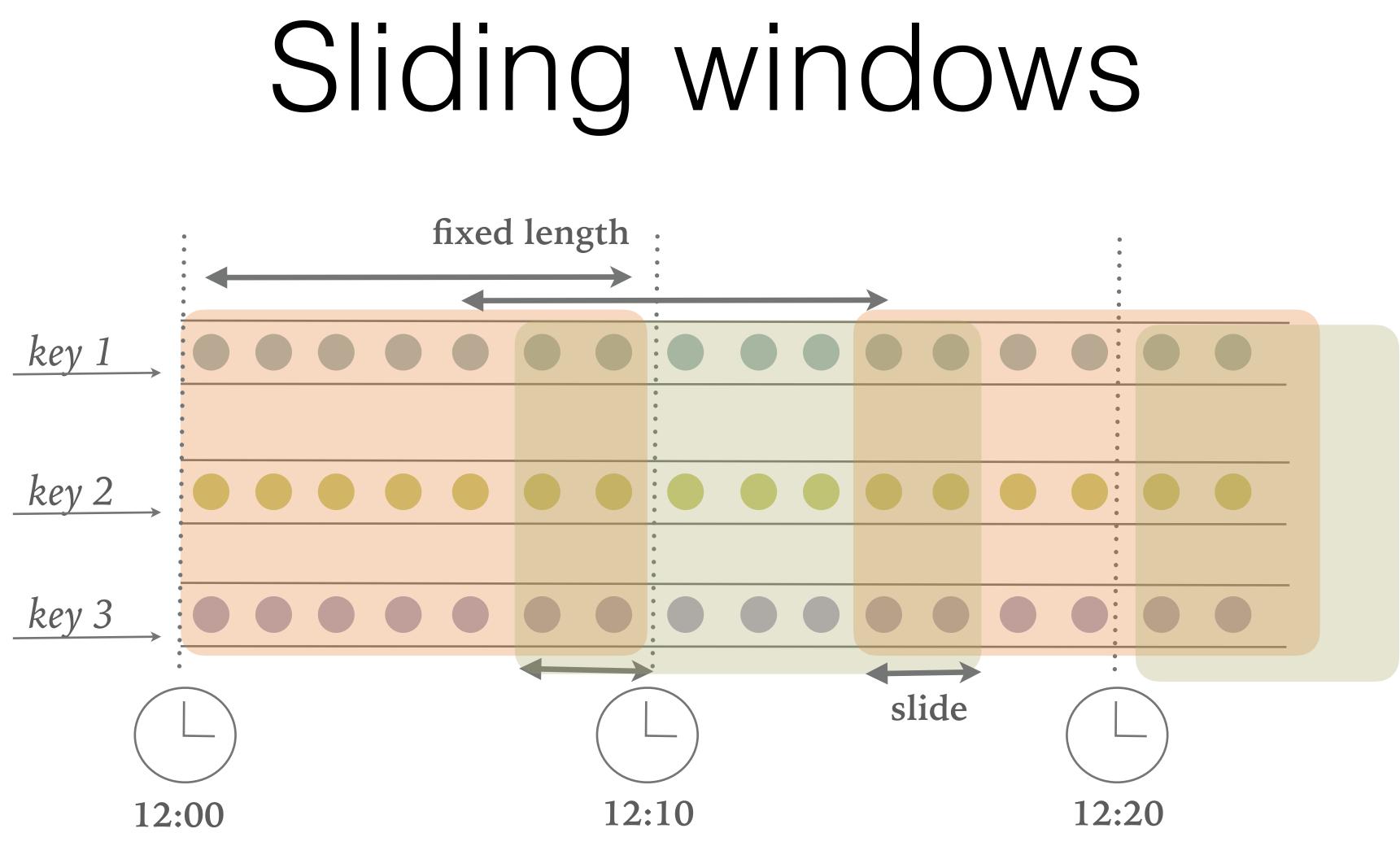


Tumbling windows



non-overlapping buckets of fixed size

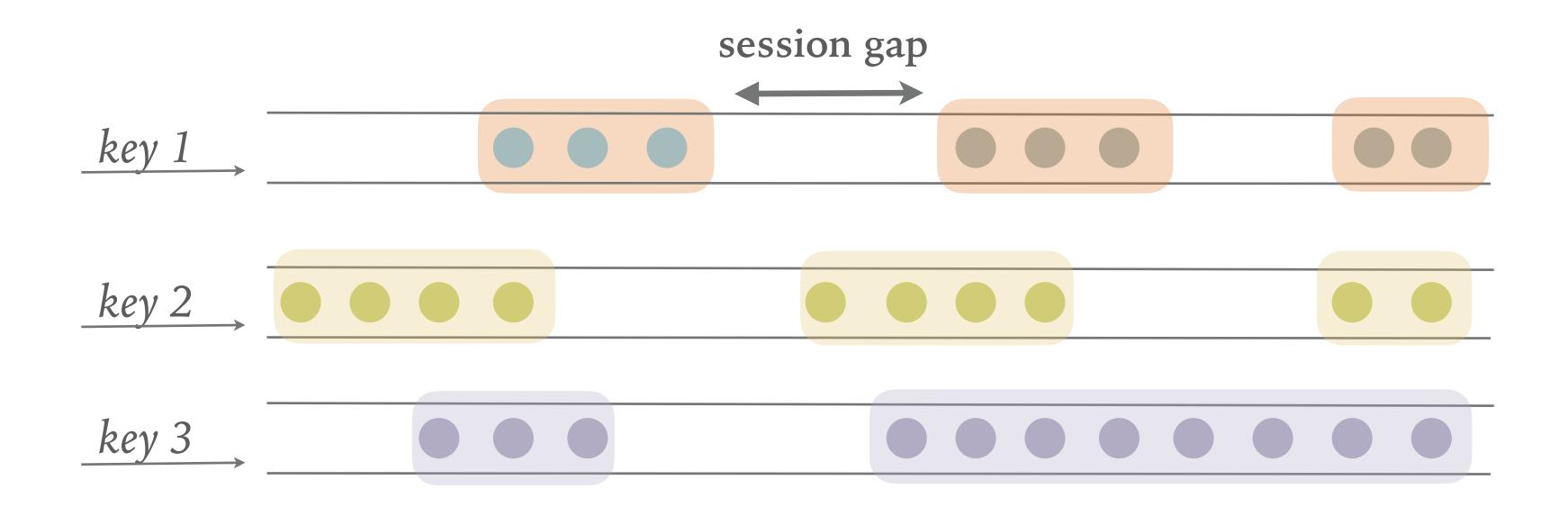




overlapping buckets of fixed size



Session windows



a period of activity followed by a period of inactivity



Windowing measures

measures:

- **Time**, such as event or processing time •
- **Count**, i.e. number of events lacksquare
- - e.g. when the amount of bids placed for an action exceeds a threshold

We can mix windowing measures to define multi-measure windows

e.g. output the last 10 tuples (count) every 5 second (time).

We can define windows based on multiple (monotonically increasing)

Data-dependent advancing measure, such as a punctuation or other signal in the stream





Window aggregation functions

Window aggregation functions define the computation we perform on the elements of a window.

- **Distributive** functions: final values can be computed as the aggregation of partial aggregates with constant size.
 - min, max, sum
- Algebraic functions: final values can be computed by applying a function on partial aggregates of fixed size.
 - average, N largest values
- Holistic functions: partial aggregates have an unbounded size
 - median, most frequent, rank



Window context

order to know where windows start and end:

- Context Free (CF) windows are those for which we can compute their boundaries without processing any tuples.
 - Tumbling and sliding windows are context-free as we can compute all start and end timestamps based on their length and slide parameters.
- Forward Context Free (FCF) windows are those which depend on punctuations. We can compute their boundaries once we have processed all events up to a timestamp t.
- Forward Context Aware (FCA) windows require us to process tuples after timestamp t in order to compute all window boundaries before t.
 - Multi-measure windows are FCA.

We divide windows into 3 classes with regard to the context we need in



Window evaluation





Window evaluation strategies

belongs to the window, or **lazily**, on trigger.

(ii) upon receiving a trigger, the operator needs to decide which windows are "complete", apply the final aggregation, and produce results.

- A window operator is responsible for grouping incoming records into windows and making the evaluation function results available in the output whenever a window triggers, i.e. when the system's notion of time arrives at its end timestamp.
- Evaluation functions can be applied **eagerly**, upon receiving a new record that
- Regardless of the strategy used, the operator performs two types of processing:
- (i) upon receiving a new event in its input, the operator needs decide how to assign the event to one or more windows and possibly apply some partial aggregation.



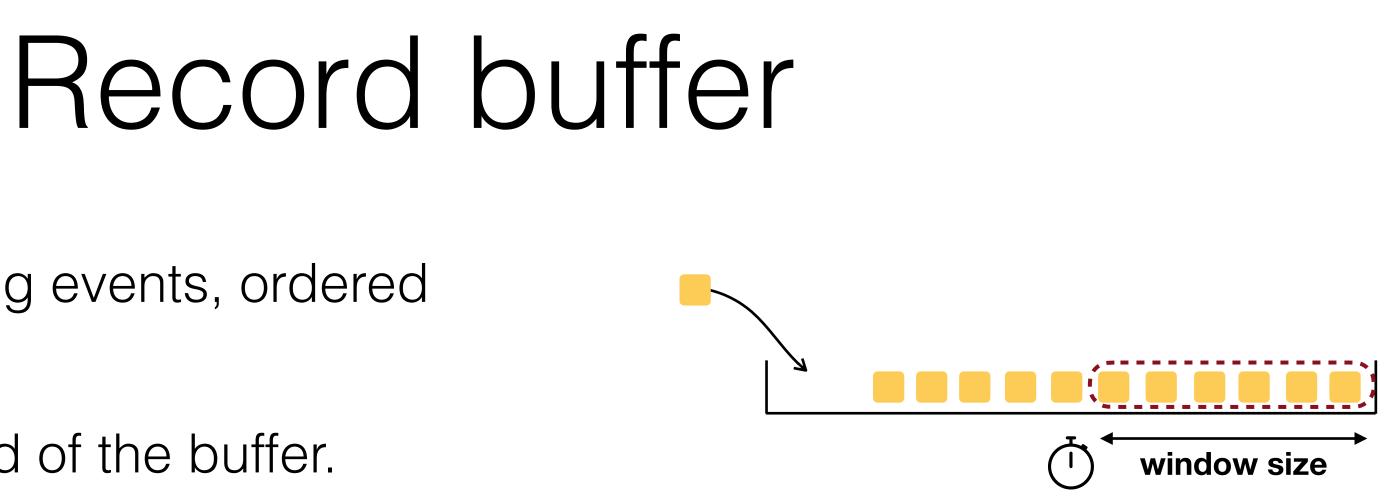
Use a buffer to store incoming events, ordered by timestamp.

On event: Append to the end of the buffer.

On trigger: Retrieve all records whose timestamp falls inside the window bounds.

The number of records in the buffer is proportional to the window size and input rate, thus, state requirements can grow significantly for high input rates and large windows.

The evaluation function is applied to the window contents lazily at trigger time.







Record buffer window size

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The evaluation function is applied to the window contents lazily at trigger time.

- General strategy which can support all window types and aggregation functions.
- Evaluation can be inefficient for \bullet out-of-order streams (memory copies), high rate (increasing state), and overlapping windows (multiple aggregate computations).



Aggregate tree

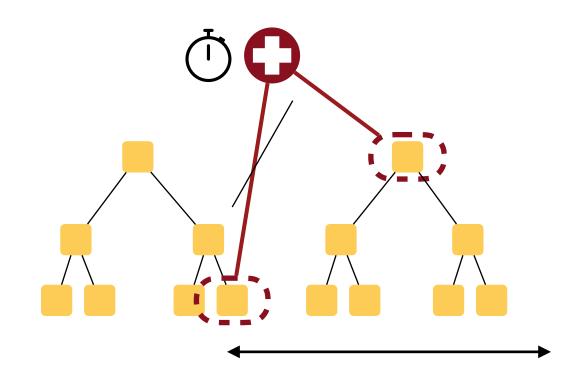
Use a binary tree to store partial aggregates on top of stream events.

On event: Insert to the binary tree and update all affected partial aggregates.

On trigger: Combine partial aggregated to compute and emit the final value.

The state requirements are high as both events and partial aggregates are maintained.

The evaluation function is applied on event and on trigger.





window size



Aggregate tree

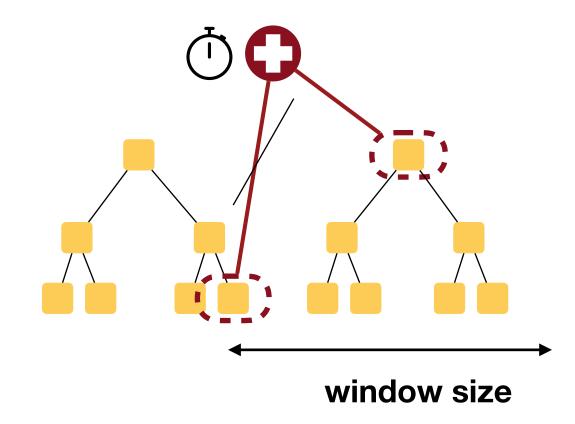
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- Partial aggregates can be shared across overlapping windows.
- Memory copy for out-of-order events or lacksquaretree re-balancing.
- Low-latency trigger as final aggregates can be computed by combining the pre-computed partial aggregates.



Organize events into windows by assigning them IDs (start or end timestamp).

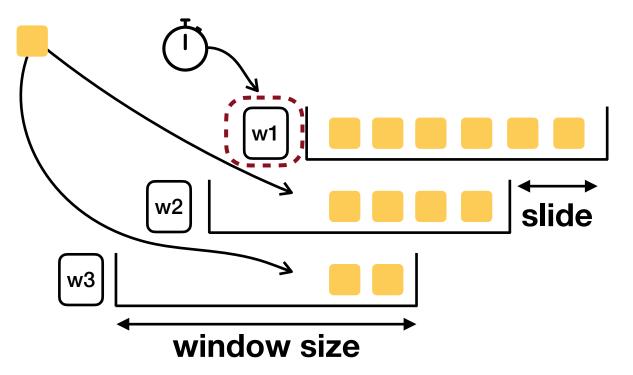
On event: An assigner function computes a list of at most $\left|\frac{1}{1}\right|$ window IDs the record belongs to. The event is inserted to each window in the list.

On trigger: Window contents are retrieved using the ID.

This strategy has low state requirements when the evaluation function is associative and commutative and can be eagerly applied on record arrival.

If the window slide is much smaller than the window length, successive windows have large overlap, resulting to redundant computations and high state requirements.

Window ID (Bucket)







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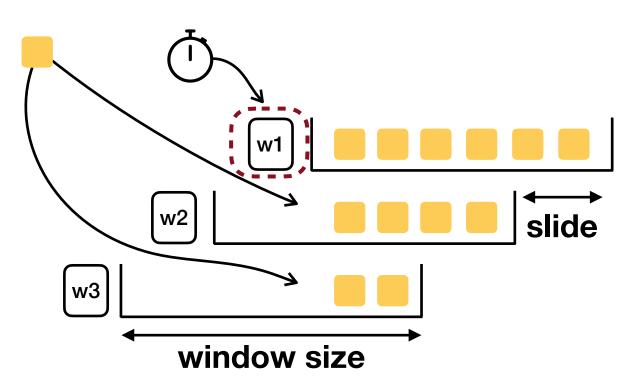
On event: An assigner function computes a list of at most $\left|\frac{1}{r}\right|$ window IDs the record belongs to. The event is inserted to each window in the list.

On trigger: Window contents are retrieved using the ID.

This strategy has low state requirements when the evaluation function is associative and commutative and can be eagerly applied on record arrival.

If the window slide is much smaller than the window length, successive windows have large overlap, resulting to redundant computations and high state requirements.

Window ID (Bucket)



- Low latency on trigger if aggregation \bullet can be computed eagerly.
- Redundancy, high memory requirements, and high latency on event (many assignments) for overlapping windows.



Window Slicing

Organize events into smaller units, called **panes**.

A pane is the maximum shareable unit across windows and its size is computed as $l_p = gcd(l, l_s)$.

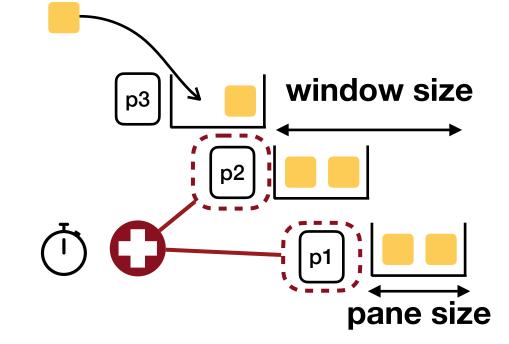
This guarantees that a record belongs to only one pane and that every window can be composed by a set of consecutive panes.

On event: An assigner function computes the pane ID and adds the record to its state.

On trigger: Retrieve $\frac{1}{7}$ panes to assemble the window contents. ι_p

If the evaluation function supports pre-aggregation, it can be eagerly applied on record arrival to maintain partially aggregated results per pane.

The final aggregate is computed by combining the partial aggregates on trigger.







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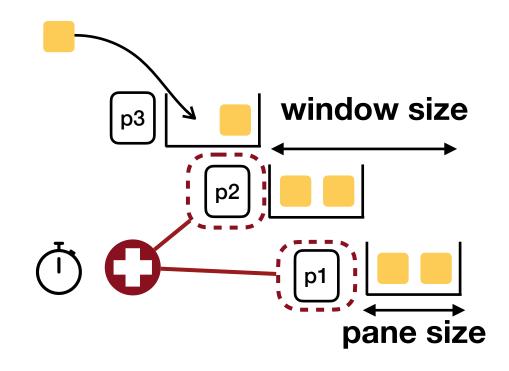
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The final aggregate is computed by combining the partial aggregates on trigger.



- Low state requirements and low \bullet latency on trigger.
- When the window length is a ulletmultiple of the window slide, results can be shared across windows.



Flink window functions





Window functions

Window functions define the computation that is performed on the elements of a window

- added to a window:
 - aggregated value as the result.
 - ReduceFunction and AggregateFunction
- the list of all collected elements when evaluated:
 - They require more space but support more complex logic.
 - ProcessWindowFunction

Incremental aggregation functions are applied when an element is

They maintain a single value as window state and eventually emit the

Full window functions collect all elements of a window and iterate over



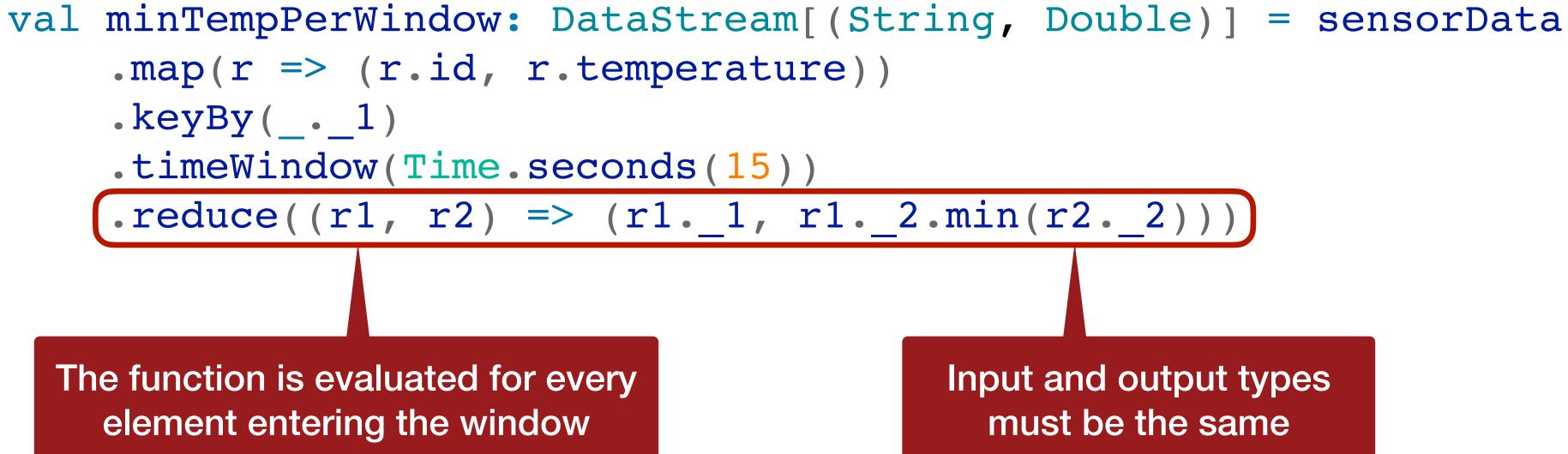
ReduceFunction example

val minTempPerWindow: DataStream[(String, Double)] = sensorData .map(r => (r.id, r.temperature))

- .keyBy(. 1)
- .timeWindow(Time.seconds(15))
- .reduce((r1, r2) => (r1. 1, r1. 2.min(r2. 2)))



ReduceFunction example





AggregateFunction interface

public interface AggregateFunction<IN, ACC, OUT> extends Function, Serializable {

// create a new accumulator to start a new aggregate. ACC createAccumulator();

ACC add(IN value, ACC accumulator);

// compute the result from the accumulator and return it. OUT getResult(ACC accumulator);

// merge two accumulators and return the result. ACC merge(ACC a, ACC b);

- // add an input element to the accumulator and return the accumulator.

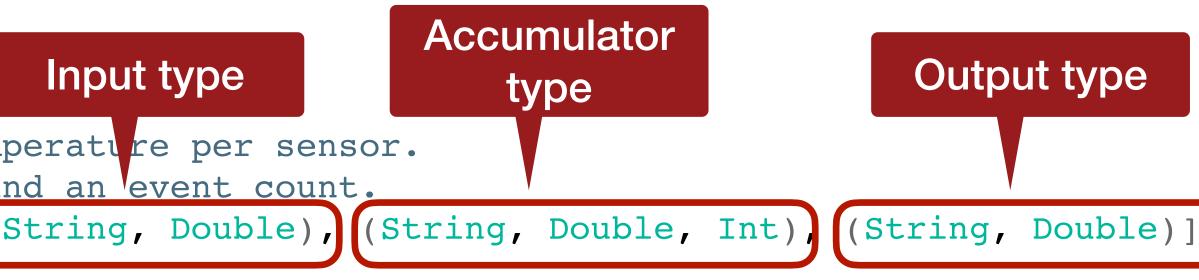


```
val avgTempPerWindow: DataStream[(String, Double)] = sensorData
    .map(r => (r.id, r.temperature))
    .keyBy( . 1)
    .timeWindow(Time.seconds(15))
    .aggregate(new AvgTempFunction)
// An AggregateFunction to compute the average temperature per sensor.
// The accumulator holds the sum of temperatures and an event count.
class AvgTempFunction extends AggregateFunction [(String, Double), (String, Double, Int), (String, Double)]
{
    override def createAccumulator() = { ("", 0.0, 0)}
    override def add(in: (String, Double), acc: (String, Double, Int)) = {
        (in._1, in._2 + acc._2, 1 + acc._3)
    override def getResult(acc: (String, Double, Int)) = { (acc._1, acc._2 / acc._3) }
    override def merge(acc1: (String, Double, Int), acc2: (String, Double, Int)) = {
```

```
(acc1. 1, acc1. 2 + acc2. 2, acc1. 3 + acc2. 3)
}
```



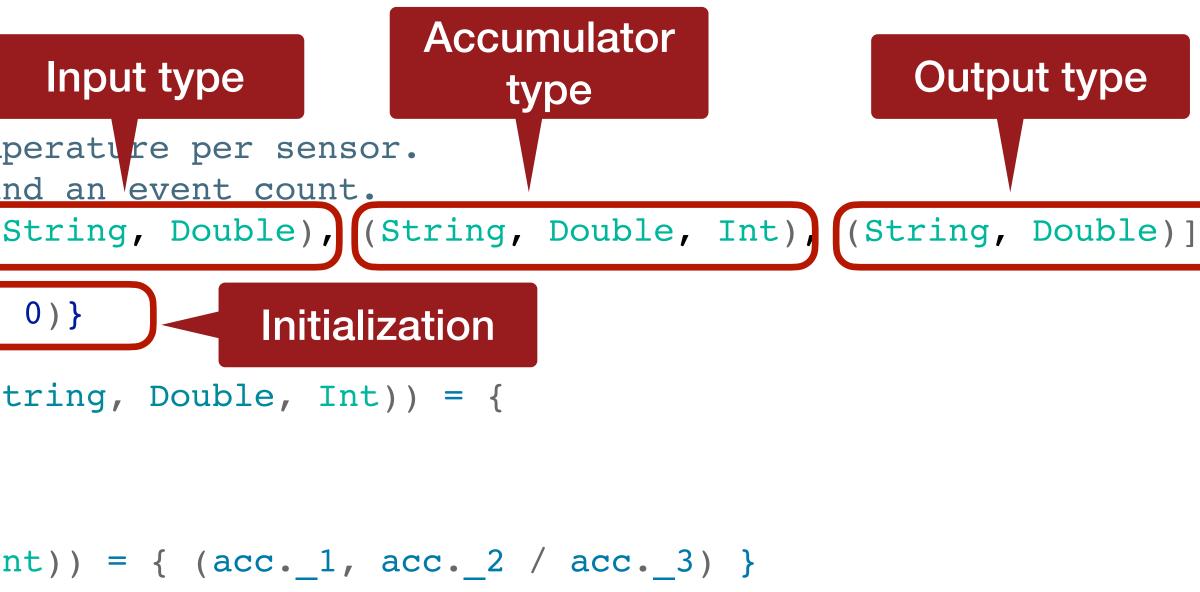
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    }
```







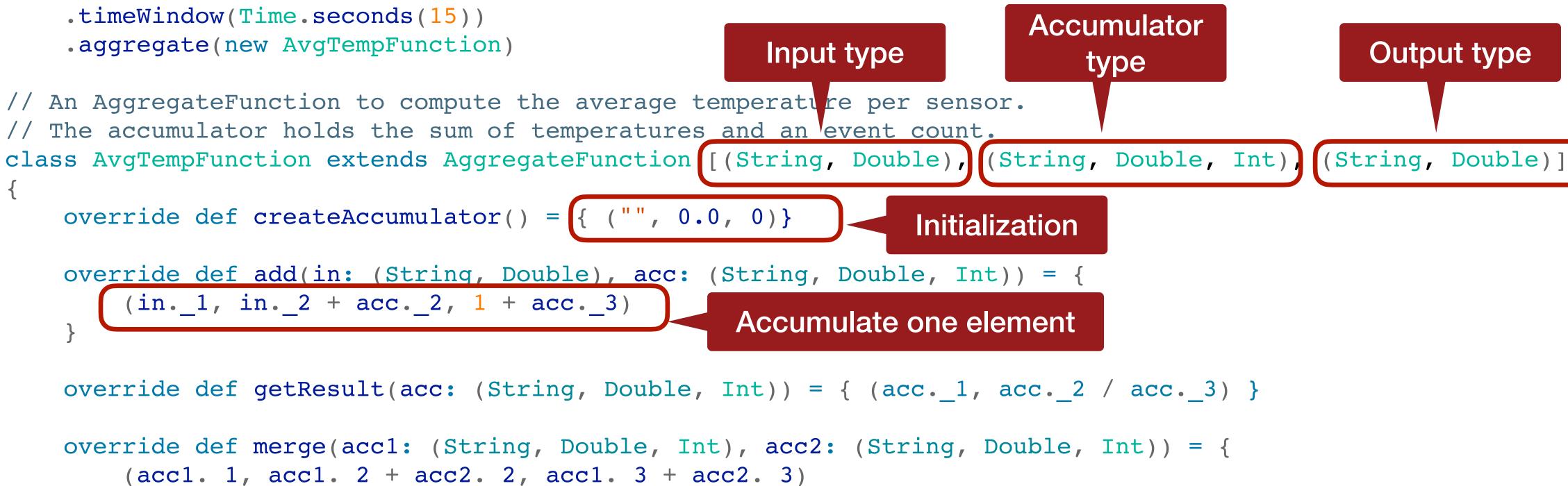
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                                                               Initialization
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```







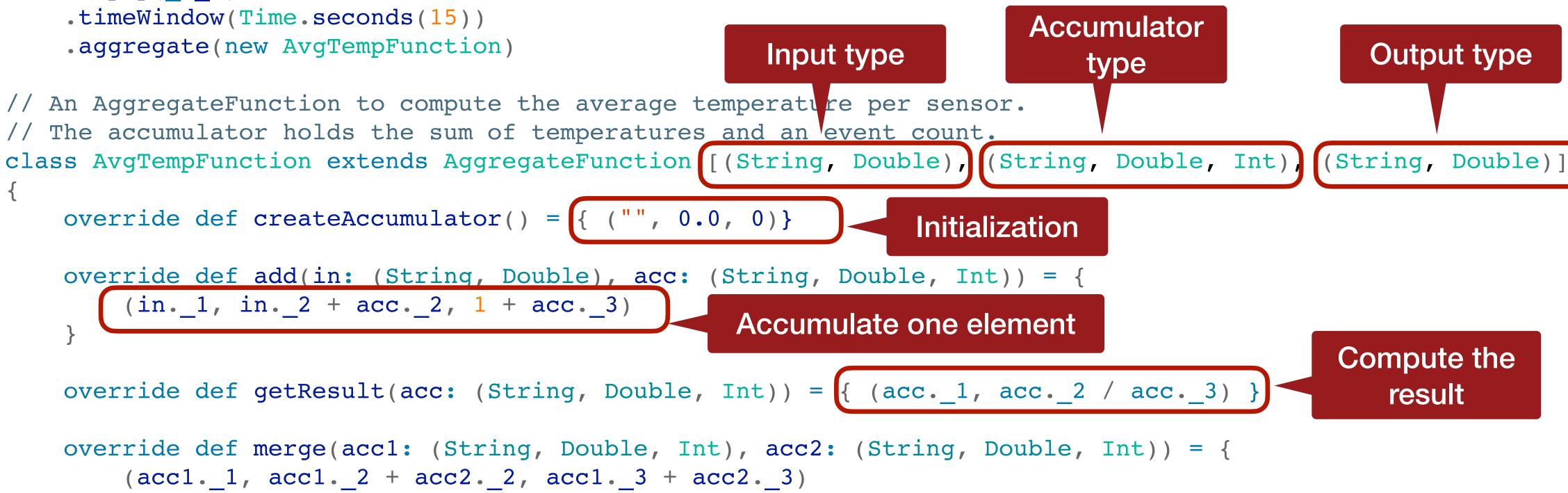
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    }
```





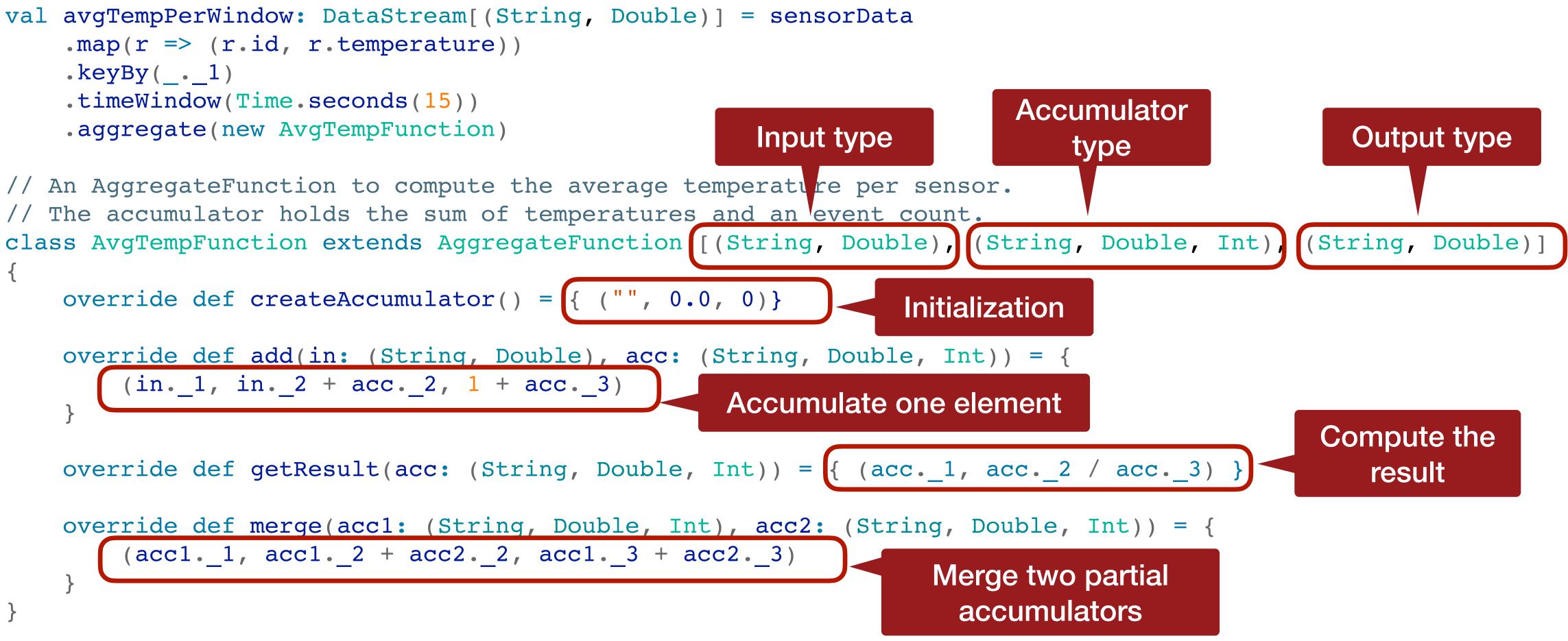


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    }
```











ProcessWindowFunction

the contents of a window:

- The process() method is called with the key of the window, an Iterable to access the elements of the window, and a Collector to emit results.
- A Context gives access to the metadata of the window (start and end timestamps in the case of a time window), the current processing time and the watermark.

Use the ProcessWindowFunction to perform arbitrary computations on



ProcessWindowFunction interface

public abstract class ProcessWindowFunction<IN, OUT, KEY, W extends Window> extends AbstractRichFunction {

// Evaluates the window void process(KEY key, Context ctx, Iterable<IN> vals, Collector<OUT> out) throws Exception;

public abstract class Context implements Serializable {

// Returns the metadata of the window public abstract W window();

// Returns the current processing time public abstract long currentProcessingTime();

// Returns the current event-time watermark public abstract long currentWatermark();



ProcessWindowFunction interface

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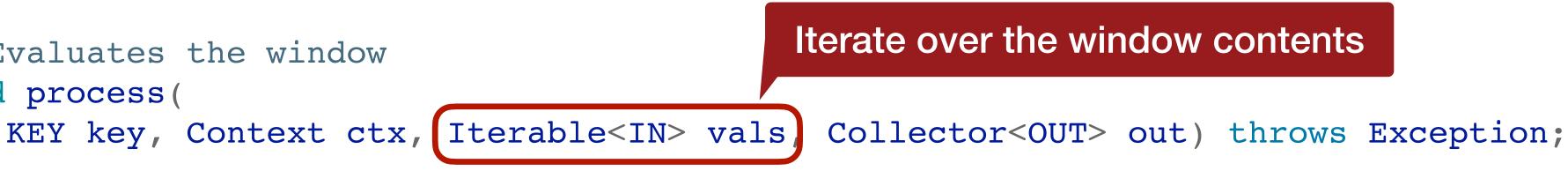
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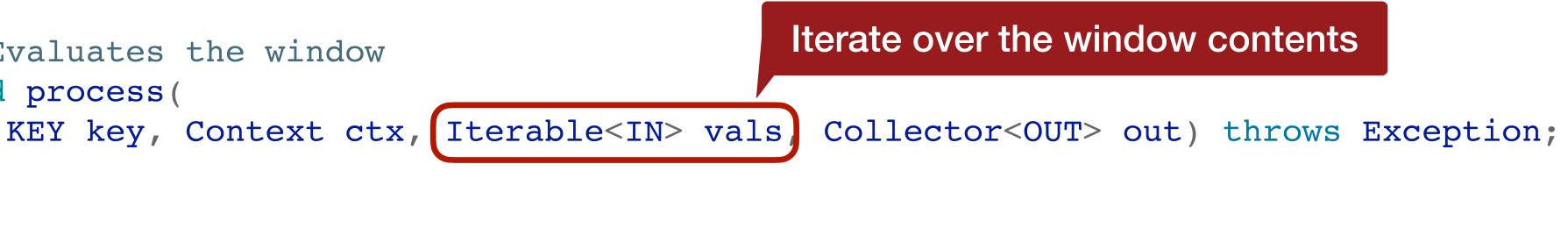
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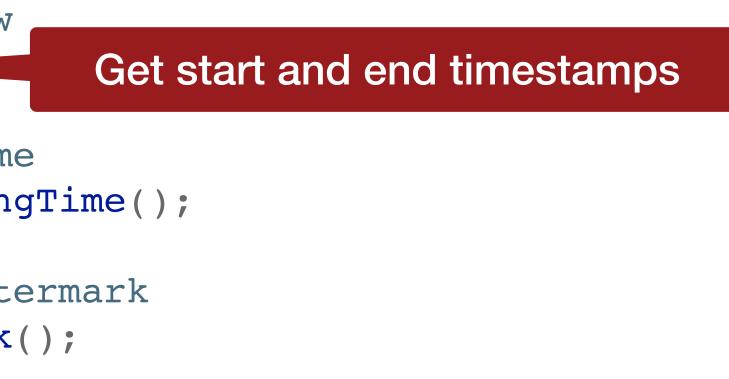
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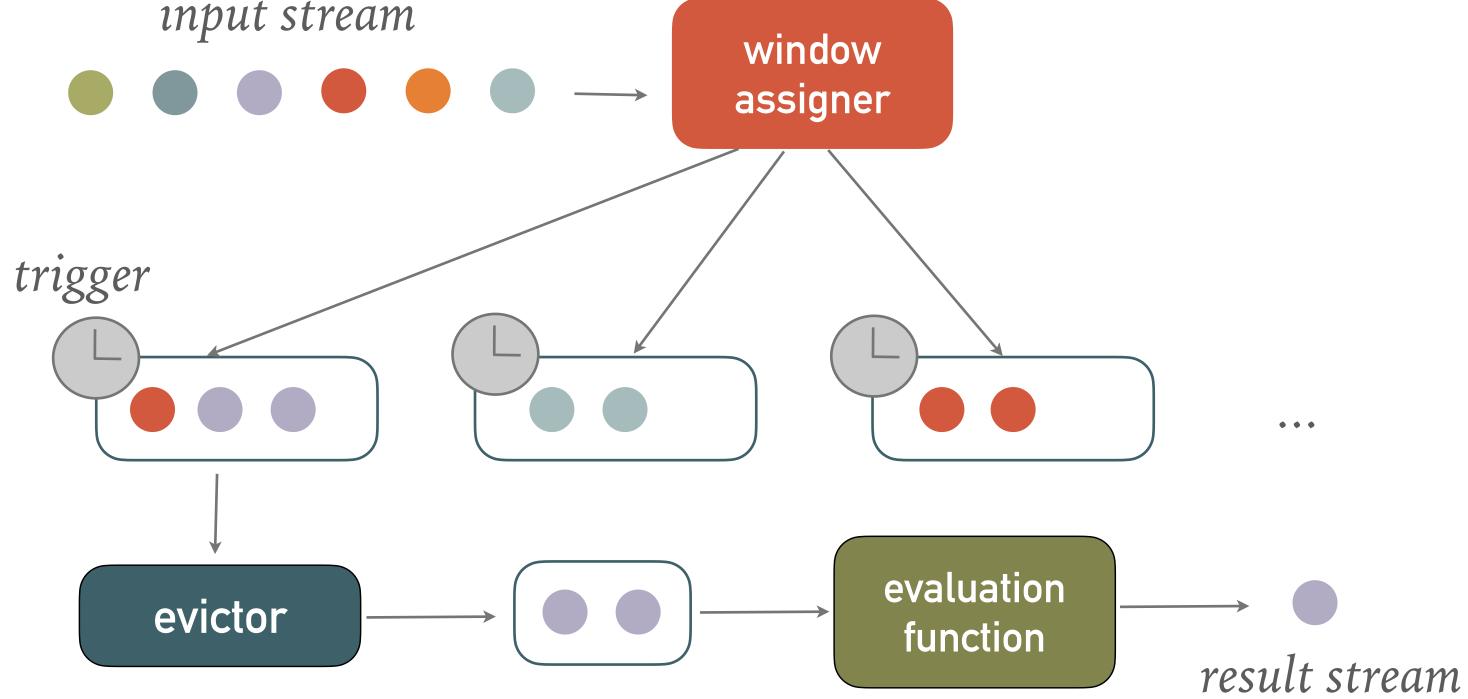






Custom windows

input stream





- Advanced transformation functions used to implement custom logic for which predefined windows and transformations might not be suitable: they provide access to record timestamps and watermarks
- they can register timers that trigger at a specific time in the future

ProcessJoinFunction, BroadcastProcessFunction, KeyedBroadcastProcessFunction, ProcessWindowFunction, and ProcessAllWindowFunction.

Process Functions

ProcessFunction, KeyedProcessFunction, CoProcessFunction,



The KeyedProcessFunction is applied to a KeyedStream:

- each record of the stream. Result records are emitted by passing them to the current record and to a TimerService.
- timestamp of the firing timer and the Collector allows emitting records. The OnTimerContext provides the same services as the Context object of the time) of the firing timer.

KeyedProcessFunction

• processElement(v: IN, ctx: Context, out: Collector[OUT]) is called for Collector. The Context object gives access to the timestamp and the key of the

• onTimer(timestamp: Long, ctx: OnTimerContext, out: Collector[OUT]) is invoked when a previously registered timer triggers. The timestamp argument gives the processElement() method and also returns the time domain (processing time or event



```
val warnings = readings
    .keyBy( .id) // key by sensor id
    .process(new TempIncreaseAlertFunction) // apply ProcessFunction to monitor temperatures
/** Emits a warning if the temperature of a sensor monotonically increases for 1 second (in processing time) **/
class TempIncreaseAlertFunction extends KeyedProcessFunction[String, SensorReading, String] {
    // stores temperature of last sensor reading
    val lastTemp: Double
    // stores timestamp of currently active timer
```

```
val currentTimer: Long
```

```
override def processElement(r: SensorReading, ctx:Context, out: Collector[String]): Unit = {
    // get previous temperature
   val prevTemp = lastTemp
    // update last temperature
   lastTemp = r.temperature
```

```
if (prevTemp == 0.0 || r.temperature < prevTemp) {</pre>
    // temperature decreased; delete current timer
    ctx.timerService().deleteProcessingTimeTimer(curTimer)
} else if (r.temperature > prevTemp && curTimerTimestamp == 0) {
    // temperature increased and we have not set a timer yet: set processing time timer for now + 1 second
    val timerTs = ctx.timerService().currentProcessingTime() + 1000
    ctx.timerService().registerProcessingTimeTimer(timerTs)
// remember current timer
currentTimer = timerTs
```



onTimer() example

```
override def onTimer(
    ts: Long,
    ctx: OnTimerContext,
    out: Collector[String]): Unit = {
        out.collect("Temperature of sensor '" + ctx.getCurrentKey +
            "' monotonically increased for 1 second.")
        currentTimer.clear()
```



CoProcess Function

For low-level operations on two inputs:

- One transformation method for each input processElement1() and processElement2()
- Both methods are called with a Context object that gives access to the element or timer timestamp and a TimerService
- You can use it to register timers and it provides an onTimer() callback method



val forwardedReadings = readings .connect(filterSwitches) // key by sensor ids .keyBy(.id, .1) .process(new ReadingFilter)

```
// connect readings and switches
// apply filtering CoProcessFunction
```

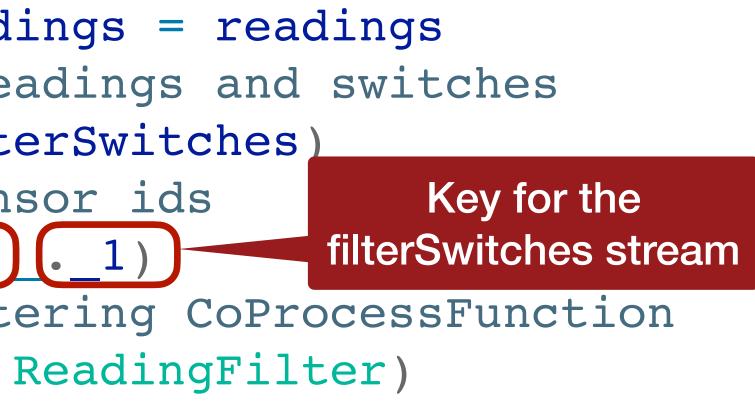


7	val forwardedReadi	ng
	// connect read	di
Key for the readings	.connect(filte:	rS
stream	// key by sense	or
	.keyBy(id,	•
	<pre>// apply filte:</pre>	ri
	.process(new Re	ea

```
s = readings
              ngs and switches
              Switches)
               ids
              1)
              .ng CoProcessFunction
rocess(new ReadingFilter)
```



7	val forwardedReading
	// connect readi
Key for the readings	.connect(filterS
stream	<pre>// key by sensor</pre>
	.keyBy(id,
	<pre>// apply filteri:</pre>
	.process(new Read





class ReadingFilter extends CoProcessFunction[SensorReading, (String, Long), SensorReading] {

// process readings override def processElement1(reading: SensorReading, ctx: Context, out: Collector[SensorReading]): Unit = {...}

// process switches override def processElement2(switch: (String, Long), ctx: Context, out: Collector[SensorReading]): Unit = {...}

// process timers override def onTimer(ts: Long, ctx: OnTimerContext, out: Collector[SensorReading]): Unit = {...}



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